

# Construction and Performance of Network Codes.

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Cerdanyola del Vallès (Barcelona), 2013.

- 1 PNC
- 2 Lattice Network Codes
- 3 C&F
- 4 C&F HAMMING
- 5 Improvement of the Coefficients
- 6 Conclusions

# Outline

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# Physical-layer Network Coding

## Introduction

### Today

Interference is treated as a destructive phenomenon.

### Network Coding Introduced the Idea

Intermediate nodes in a network are able to perform operations to the input packets rather than just forwarding them.

### Network Coding at the Physical-layer? PNC

When multiple electromagnetic waves come together within the same physical space, they add. This additive mixing of electromagnetic waves is a form of Network Coding, performed by nature. PNC aims to exploit this fact.

# Physical-layer Network Coding

## Main ideas

### The Source Transmits a Message

$w_l \in \mathbb{F}_p^k$ , where  $\mathbb{F}_p$  is a finite field with  $p$  elements  $\{0, 1, 2, \dots, p-1\}$  and  $p$  is a prime number.

### The Relay Decodes a Linear Combination $v$ of these Messages

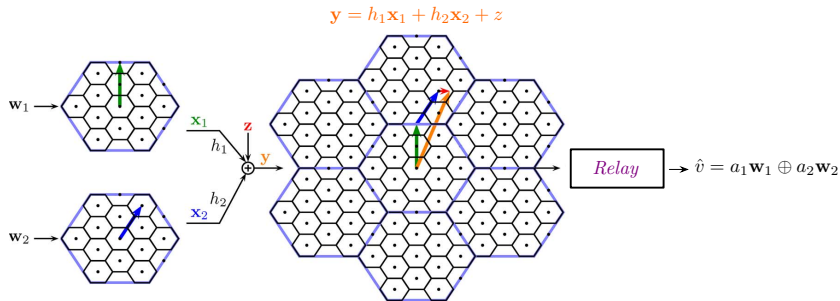
$v = a_1 w_1 \oplus a_2 w_2 \oplus \dots \oplus a_L w_L$ , where  $a_l$  are coefficients over the finite field  $\mathbb{F}_p$ .

### The Destination Can Solve For the Original Messages if

$A = \begin{bmatrix} a_{11} & a_{12} & a_{1L} \\ a_{21} & a_{22} & a_{2L} \\ \vdots & \vdots & \vdots \\ a_{M1} & a_{M2} & a_{ML} \end{bmatrix}$  has rank  $L$ .

# Physical-layer Network Coding

## Example



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# Lattice Network Codes

## Introduction

### What Are We Looking For?

If the waveforms at the transmitter are points of a lattice (that is  $\mathbb{Z}$  or  $\mathbb{Z}[i]$ ), then every integer combination of these waveforms is itself a point of the same lattice.

### The Algebraic Structure Necessary Is

Given a  $R$ -lattice  $\Lambda$  (e.g.  $\mathbb{Z}[i]$ ) and a sublattice  $\Lambda'$  of  $\Lambda$  (e.g.  $\pi\mathbb{Z}[i]$ ), the quotient group  $\Lambda/\Lambda'$  (e.g.  $\frac{\mathbb{Z}[i]}{\pi\mathbb{Z}[i]}$ ) is a  $R$ -module. For a Lattice Network Code, the message space is  $W = \Lambda/\Lambda'$ .

### Let's See a Bit More of Insight

The  $R/aR$  structure, being  $R$  a PID and  $a$  prime, forms a field. Thus, we will be able to find an isomorphism between  $\mathbb{F}_p = \frac{\mathbb{Z}}{p\mathbb{Z}}$  and  $\frac{\mathbb{Z}[i]}{\pi\mathbb{Z}[i]}$  if both fields have the same number of elements.



# Lattice Network Codes

## The Lattice $\mathbb{Z}[i]$

### GAUSSIAN Integers

$$\mathbb{Z}[i] = \{a + bi \mid a, b \in \mathbb{Z}\}.$$

### Prime Factorization Used

If  $p \equiv 1 \pmod{4}$  then  $p = \pi\pi^*$  is a product of two conjugate primes  $\pi, \pi^*$ .

### Example

The prime  $p = 5$  satisfies  $5 \equiv 1 \pmod{4}$ , so 5 has two conjugate GAUSSIAN prime factors.

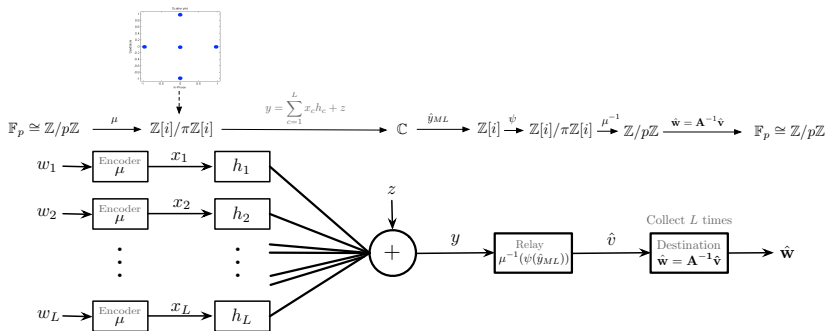
$$\text{Since } 5 = 1^2 + 2^2, 5 = (1 + 2i)(1 - 2i).$$

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- 3 C&F
  - C&F System: Scalar Case
  - C&F System: Vectorial Case

# C&F System: Scalar Case

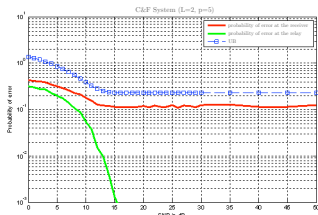
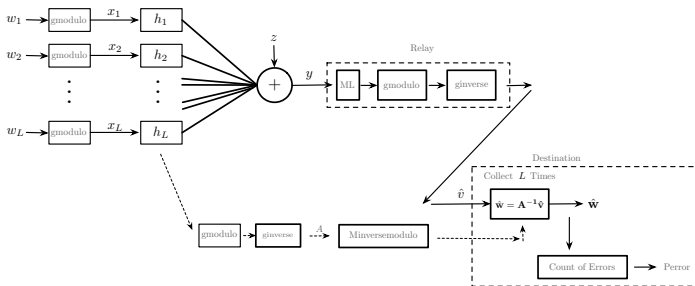
## System Model



$$\begin{aligned} \mu(w_i) &= w_i \bmod \pi = w_i - \left\lfloor \frac{w_i \pi^*}{\pi \pi^*} \right\rfloor \pi \\ \mu^{-1}(z) &= z \bmod p = (z^* u \pi + z v \pi^*) \bmod p \\ y &= h_1 x_1 + h_2 x_2 + \dots + h_L x_L + z \\ \hat{v} &= a_1 w_1 \oplus a_2 w_2 \oplus \dots \oplus a_L w_L \end{aligned}$$

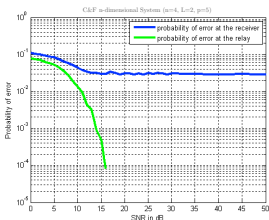
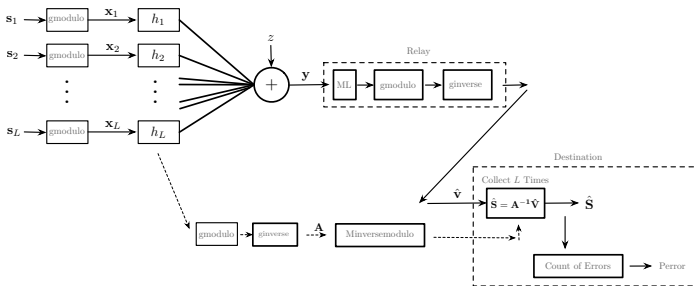
# C&F System: Scalar Case

## Performance



- 3** C&F
  - C&F System: Scalar Case
  - C&F System: Vectorial Case

# C&F System: Vectorial Case Performance



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# HAMMING $q$ -ary Codes

## Generating Matrix

$$C = uG : u \in \mathbb{F}_p^k.$$

We say that  $G$  is systematic if  $G = (I_k | -P^T)$ .

## Parity Check Matrix

$$C = \{v \in \mathbb{F}_p^n : Hv^T = 0\}.$$

If  $G$  is systematic, a parity check matrix is  $H = (P | I_{n-k})$ .

# C&F HAMMING $q$ -ary Coded System Construction

## HAMMING $q$ -ary Code

Given an integer  $r \geq 2$ , let  $n = \frac{q^r-1}{q-1}$ . The HAMMING  $q$ -ary code is a linear  $[n, n-r]$  code in  $\mathbb{F}_q^n$ , whose parity check matrix  $H$  is such that

$$H = (v_1 | v_2 | \dots | v_n)$$

where  $v_1, \dots, v_n \in \mathbb{F}_q^r$  is a list of nonzero vectors satisfying the condition that no two vectors are scalar multiples of each other.

## Example

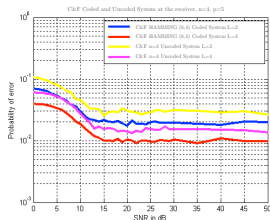
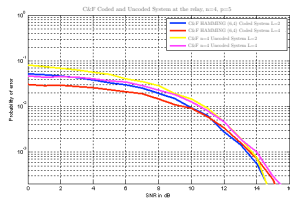
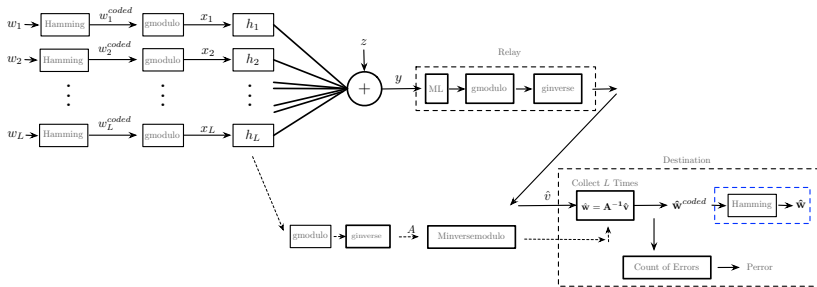
Let  $\mathbb{F}_5$  and  $r = 2$ ,  $n = \frac{5^2-1}{5-1} = 6$ . So,  $k = n - r = 4$ . A straightforward way to generate a systematic HAMMING  $q$ -ary code is generating the matrix  $P$  as a  $r \times k$  matrix with columns

$$P = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 2 & 3 & 4 \end{bmatrix}.$$

And then generate  $H$  and  $G$  using  $G = (I_k | -P^T)$  and  $H = (P | I_{n-k})$ .

$$H = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 0 \\ 1 & 2 & 3 & 4 & 0 & 1 \end{bmatrix} \quad \text{and} \quad G = \begin{bmatrix} 1 & 0 & 0 & 0 & 4 & 4 \\ 0 & 1 & 0 & 0 & 4 & 3 \\ 0 & 0 & 1 & 0 & 4 & 2 \\ 0 & 0 & 0 & 1 & 4 & 1 \end{bmatrix}.$$

# C&F HAMMING $q$ -ary Coded System Performance

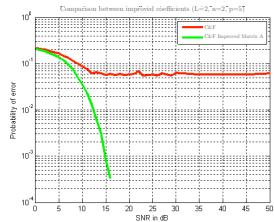
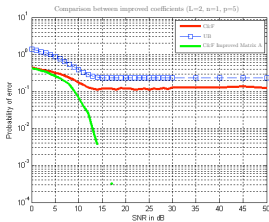
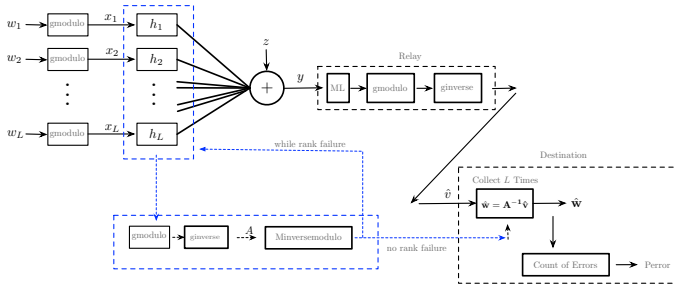


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## 5 Improvement of the Coefficients

- Improvement of the Coefficients: Improved Matrix A
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# Improved Matrix A Performance



## 5 Improvement of the Coefficients

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# Optimum Matrix A Construction

## Scalar Factor

$$\beta_{MMSE} = \frac{\text{SNR} h^T a_m}{\text{SNR} \|h_m\|^2 + 1}.$$

## Optimum Coefficients

**Theorem:** For a given vector of coefficients of the channel  $h_m = [h_{m1}, h_{m2}, \dots, h_{mL}]^T \in \mathbb{R}^L$ , the computation rate is maximized by choosing in network coding the vector of coefficients  $a_m \in \mathbb{Z}^L$  as

$$a_m = \arg \min_{a_m \in \mathbb{Z}^L, a_m \neq 0} (a_m^T G_m a_m)$$

where

$$G_m = I - \frac{\text{SNR}}{1 + \text{SNR} \|h_m\|^2} H_m.$$



# Optimum Matrix A Construction

## What's Behind this Minimization?

$$a_m = \arg \min_{a_m \in \mathbb{Z}^L, a_m \neq 0} (a_m^T G_m a_m).$$

- CHOLESKY factorization.
- Lattice reduction: LLL algorithm.
- Vector search: SCHNORR EUCHNER method.

# Optimum Matrix A

## Solving the ILS Problem

### ILS Problem

$$\min_{z \in \mathbb{Z}^n} \|y - Bz\|^2$$

this problem is analogous to solving

$$\min_{z \in \mathbb{Z}^n} (y - Bz)^T V^{-1} (y - Bz).$$

One can first compute the CHOLESKY factorization  $V = R^T R$ , then solve two lower triangular linear systems  $R^T \bar{y} = y$  and  $R^T \bar{B} = B$ .

As our real aim is to solve the SVP problem

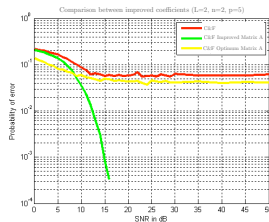
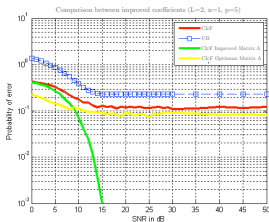
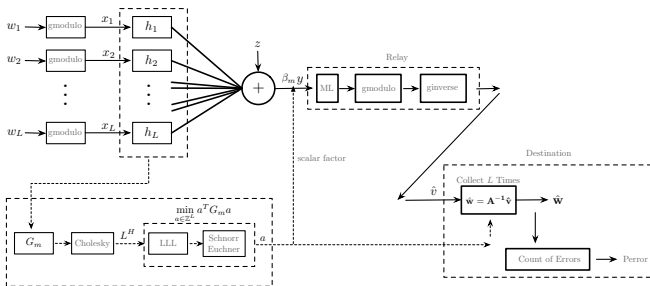
$$\min_{z \in \mathbb{Z}^n} (z)^T V^{-1} (z)$$

we use  $B = -I_n$  and  $y = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix}_n$  and therefore  $\bar{B} = R^T B$  and  $\bar{y} = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix}_n$ .

Finally the problem becomes

$$\min_{z \in \mathbb{Z}^n} \|\bar{y} - \bar{B}z\|^2.$$

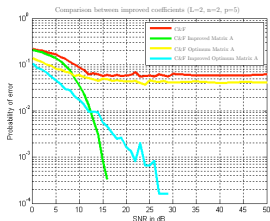
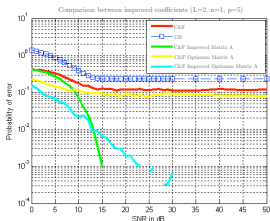
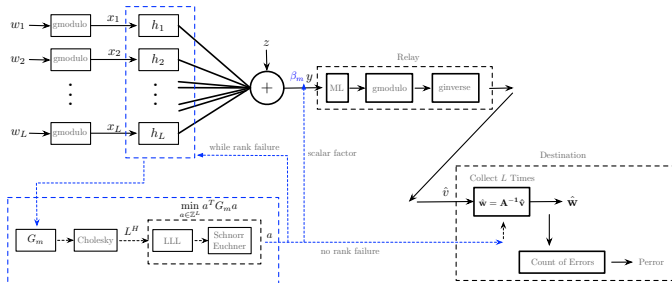
# Optimum Matrix A Performance



## 5 Improvement of the Coefficients

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# Improved Optimum Matrix A Performance



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# Conclusions

- Mathematical tools. ✓
- C&F uncoded system model scalar:
  - MATLAB  $L$ -dimensional  $\forall p$  implementation using a working  $L = 2, p = 5$  code base with given ML detector. ✓
- C&F uncoded system model vectorial:
  - MATLAB implementation  $n$ -dimensional. ✓
- C&F HAMMING  $q$ -ary coded system model:
  - MATLAB implementation C&F HAMMING (6,4) coded system  $n = 4$ . ✓
- Improvement of the Coefficients:
  - MATLAB implementation improved matrix  $A$ . ✓
  - MATLAB implementation optimum matrix  $A$ . ✓
  - MATLAB implementation improved optimum matrix  $A$ . ✓
- Implementation of sphere decoder for ML detection:
  - Adapting the code used for optimum matrix  $A$  as an efficient sphere decoder. ✓

## Obtained Results

- We have explained the lattice theory needed.
- We have provided several MATLAB code implementations for C&F system.
- The results of the improvement of the coefficients show:
  - Improved optimum matrix  $A$  works really well for SNR low.
  - Improved matrix  $A$  has a better slope performance for SNR high.



Thank You

**UAB**

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DE CURTÓ I DÍAZ Joaquim.

Thank you.